## Test $\mathbf{2}$

## Algorithms

- 1. (15) Consider the union-find problem on a set of n elements  $\{1, 2, ..., n\}$ . Initially each element is a subset by itself. There are two types of operations on these subsets of S.
  - (a) The find operation F(x) returns the subset to which the element x belongs.
  - (b) The union operation U(u, v) makes the two subsets u and v into one subset.

Note that each subset has a unique name, and the name cannot be changed. Suppose that each subset is represented by a tree. Each vertex has a pointer points to its parent, except the root whose pointer points to itself.

- (a) (5) Assume that n = 7, and each subset is named by the smallest element in that subset. Show the subsets by drawing the forest after each of the following sequence of union operations: U(1,2), U(3,4), U(1,3), U(5,6), U(1,5).
- (b) (10) Design efficient algorithms for the union and find operations so that each find operation can be done in  $O(\log n)$  time and the union can done in O(1) time.
- 2. (15) A bipartite graph is a graph whose vertex set V can be partitioned into two subsets  $V_1$  and  $V_2$ , such that no edges connecting vertices in the same partition. A k-cube is a graph whose vertex set is the set of binary strings of length k. Two vertexes u and v are connected if they differs in only 1 bit.
  - (a) (5) Draw *n*-cube for  $0 < n \leq 3$ .
  - (b) (10) Show that k-cubes are bipartite graphs for any k > 0.
- 3. (20) An Euler cycle is a closed walk which traverses each edge exactly once. An Euler path is a walk which traverses each edge exactly once. Let G be a connected undirected graph.
  - (a) Show that G has an Euler cycle if and only if the degree of every vertex is even.
  - (b) Design a linear time algorithm for finding an Euler cycle for the graph G.
- 4. (10) Given a weighted graph G = (V, E, w). Suppose that we want to find a spanning tree with maximum weight, instead of a minimum one. Design an algorithm for finding a maximum spanning tree, and justify that your algorithm is correct.
- 5. (20) Let G = (V, E, w) be a connected weighted graph. Suppose that Dijkstra's algorithm is used to solve the single source shortest path problem on the weighted graph G.
  - (a) Give an example to show that Dijkstra's algorithm does not work for graphs with negative weights.
  - (b) Can we first add a constant value to the weight of each edge to make all weights non-negative, then run the Dijkstra's algorithm to compute shortest paths? Justify your answer.
- 6. (10) Let G = (V, E, w) be a connected weighted undirected graph. Given a vertex  $s \in V$  and a shortest path tree  $T_s$  with respect to the source s, design a linear time algorithm for checking whether the shortest path tree  $T_s$  is correct or not.
- 7. (20) Let n be a product of two large primes p and q. Design a polynomial time nondeterministic algorithm for finding a factor of n, that is, p or q.
- 8. (20) The  $\{0,1\}$ -knapsack problem is defined as:

Given a set of  $K = \{(w_i, v_i)\}, 1 \le i \le n$ , and a capacity of the knapsack S, find an n-dimensional vector  $U = (u_1, u_2, \ldots, u_n)$  in  $\{0, 1\}^n$  such that  $W = \sum_{i=1}^n u_i \cdot w_i \le S$  and the value of  $P = \sum_{i=1}^n u_i \cdot v_i$  is maximized.

The *subset sum problem* is defined as:

Given a sequence of n integers  $X = x_1, x_2, \ldots, x_n$ , and an integer S. Fine a subsequence X' of X such that the sum of the integers in X' is S, or report that there are no such subsequences.

Show that the subset sum problem can be reduced to the  $\{0, 1\}$ -knapsack problem.